

Generative Operational Semantics for Relaxed Memory Models

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An operational semantics for concurrency

- Message passing concurrency well understood
 - operational/denotational models, equivalence/order relations sound type systems, proof systems, etc
- Shared-memory concurrency well understood, assuming
 - sequentially consistent execution, or
 - data race free programs
- Relaxed models used in practice
 - Compiler flexibility (source, JIT, instruction decoder/scheduler)
 - Efficiency, lock free algorithms
- Relaxed models not well understood
 - Goal: novel type system for relaxed model
 - This paper: operational semantics for soundness proof

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Transformations that occur in relaxed models

- Non-conflict reordering (conflict = same location + write)

```
p.f=0  
p.g=1 → p.g=1  
p.g=1 → p.f=0
```

- Redundant read elimination

```
x=p.f  
p.g=1  
y=p.f  
return y  
→  
x=p.f  
p.g=1  
return x
```

- Roach motel

```
x=p.f  
k.acquire() → k.acquire()  
k.acquire() → x=p.f
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Concurrency and program transformation

- Transformation correct: **no new behavior**
- Expect `p.h` incremented at most once (Dijkstra, 1965)

write g, read f

`p.g = 1`

`x = p.f`

`if(x == 0)`

`p.h ++`

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`p.h=2`

1 Background

- Sequential Consistency
- Data Race Free Model
- Java Memory Model

2 Speculative semantics

- Empirical and speculative actions
- Desirable executions allowed
- Undesirable executions prevented

3 Summary of results

- Relation to Java Memory Model
- Simulation precongruence

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Ops appear to execute in some sequential order
Ops of individual threads appear in program order
(Lamport 1977)

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          p.f = 0           x = p.f  
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Memory:

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Threads:  p.f = 0           x = p.f  
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return (1, 1) possible  
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Caching model

Indirection between action and memory

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Pending actions:

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	<code>return</code>	<code>return(x, y)</code>

Memory: `p.f=1`

Pending actions: `p.f=0`

Threads:	<code>p.f = 0</code>	<code>x = p.f</code>
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`return(1, 0)` possible

The execution has a **data race**: conflicting ops not totally ordered

Caching model

Indirection between action and memory

Program:	<i>write f twice</i>	<i>read f twice</i>
	<code>p.f = 0</code>	<code>x = p.f</code>
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Threads:	<code>p.f = 0</code>	<code>1 = p.f</code>
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`return(1, 0)` possible

The execution has a **data race**: conflicting ops not totally ordered

Data-Race Free (DRF) Semantics

- DRF programs: SC execution
- Programs with races: no comment
- Ok for C++ (Boehm and Adve, 2008)
 - no benign races
 - no safety guarantees

Java

- Defines semantics for programs with races (type safety)

- **Defined** (Gosling, Joy, Steele, 1996)
Caching semantics with “prescient reads”

- **Criticized** (Pugh 1999)
 - invalidates redundant read elimination

x=p.f	→	x=p.f
y=q.f	✗	y=q.f
z=p.f		z=x

- invalidates non-conflict reordering!

x=p.g	→	y=p.f
y=p.f	✗	x=p.g
z=q.f		z=q.f
p.x=1		p.x=1

- **Replaced** by JMM (JSR 133, 2004)
Semantics based on series of executions,
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<code>x=p.g</code>		<code>y=p.f</code>
<code>y=p.f</code>	\rightarrow	<code>x=p.g</code>
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Prescient read: seeing the future

- Caching not enough
- Program A:

<i>copy f to g</i>	<i>read g, write f</i>
x = p.f	y = p.g
p.g = x	p.f = 1
return x	return y

- In SC semantics, **return 1** impossible
- Can result from non-conflict reordering

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Threads:	x = p.f	p.f = 1
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Memory:	p.f=1	p.g=1
Threads:	1 = p.f p.g = 1 return 1	p.f = 1 <u>1 = p.g</u> return 1

Thin-air read: making things up

- Need to be careful
- Program B:

<i>copy f to g</i>	<i>copy g to f</i>
x = p.f	y = p.g
p.g = x	p.f = y
return x	return y

- **return 1** undesirable — out of “thin air”

Memory:

Threads:

x = p.f	y = p.g
p.g = x	p.f = y
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Java Memory Model (JMM)

■ Program A:

copy f to g

x = p.f

p.g = x

return x

read g, write f

y = p.g

p.f = 1

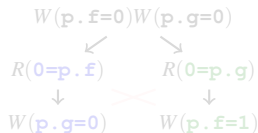
return y

■ JMM allows both threads to return 1

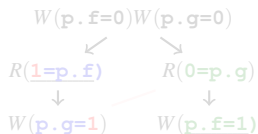
■ Criticism of JMM:

- Language acceptor, not generator
- Difficult to understand (still)
- Invalidates useful optimizations (still)
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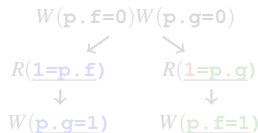
Execution 1:



Execution 2:



Execution 3:



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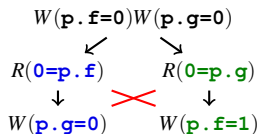
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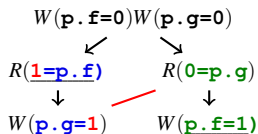
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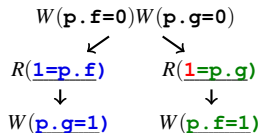
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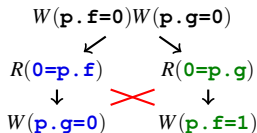
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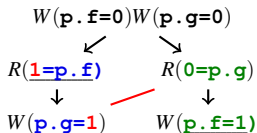
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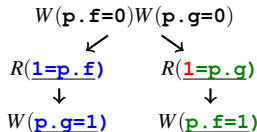
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Our goals

- New formalization of JMM
- Generative model
- Standard guarantees
 - DRF — DRF programs have SC executions
 - no TAR — no “Thin Air Reads”
- Strictly more expressive than JMM
 - Every outcome allowed by JMM allowed by our semantics
 - Only for lockless programs

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Related Work

- *The Java Memory Model*
Manson (PhD Thesis 2004)
Also Manson, Pugh, Adve, (POPL 2005)
- *Foundations of the C++ Concurrency Memory Model*
Boehm and Adve (PLDI 2008)
- *Program Transformations in Weak Memory Models*
Sevcík (PhD Thesis, 2008)
Also Sevcík and Aspinall (ECOOP 2008)
- *The semantics of x86-CC multiprocessor machine code*
Sarkar, Sewell, Nardelli, Owens, Ridge, Braibant, Myreen,
Alglave (POPL 2009)
- *Relaxed memory models: an operational approach*
Boudol and Petri (POPL 2009)
Also Boudol and Petri (ESOP 2010)

1 Background

- Sequential Consistency
- Data Race Free Model
- Java Memory Model

2 Speculative semantics

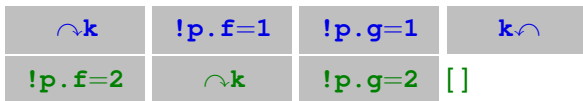
- Empirical and speculative actions
- Desirable executions allowed
- Undesirable executions prevented

3 Summary of results

- Relation to Java Memory Model
- Simulation precongruence

Memory as action sequence

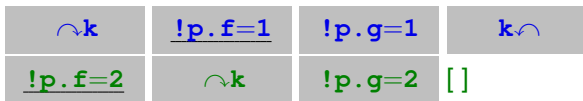
- No conventional memory
- Write action: $!p.f=1$
Acquire action: $\curvearrowright k$
Release action: $k\curvearrowleft$
- Read value determined by context (Boudol and Petri, 2009)
Context is a sequence of actions



- **Visibility** standard from JMM $p.f$ visible at 1 and 2
 $p.g$ visible at 2 only, due to k
- Threads may **reorder** non-conflict actions privately (see paper)

Memory as action sequence

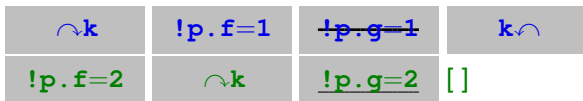
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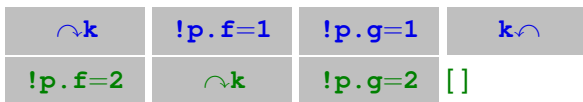
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Speculative action

- Speculation $?p . f=1$ causes branching
- Worlds execute independently: **initial** and **final**
 - Speculation visible in final branch, not initial
 - Initial branch must produce justifying empirical write $!p . f=1$

Initiality not too restrictive: Program A

- Initial branch must **justify** speculation
- Afterwards, only final copy remains

copy f to g
x = p.f
p.g = x
return x

read g, write f
y = p.f
p.g = 1
return y

~~?p.f=1~~

?p.f=1

x = p.f
p.g = x
return x
initial

y = p.g
p.f = 1
return y

x = p.f
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y = p.g
p.f = 1
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final

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~~?p.f=1~~

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p.g = 0
return 0
initial

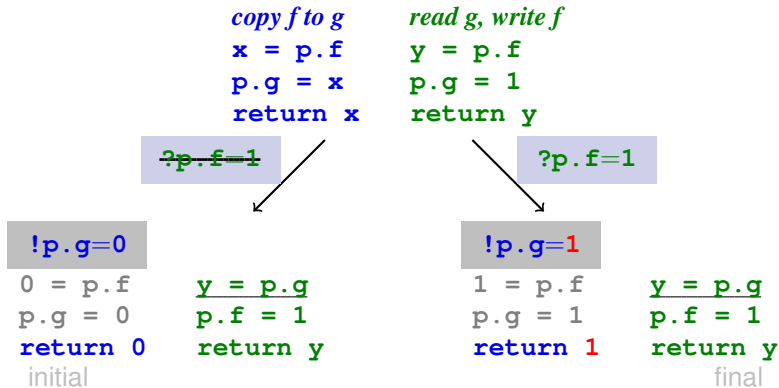
y = p.g
p.f = 1
return y

1 = p.f
p.g = 1
return 1

y = p.g
p.f = 1
return y
final

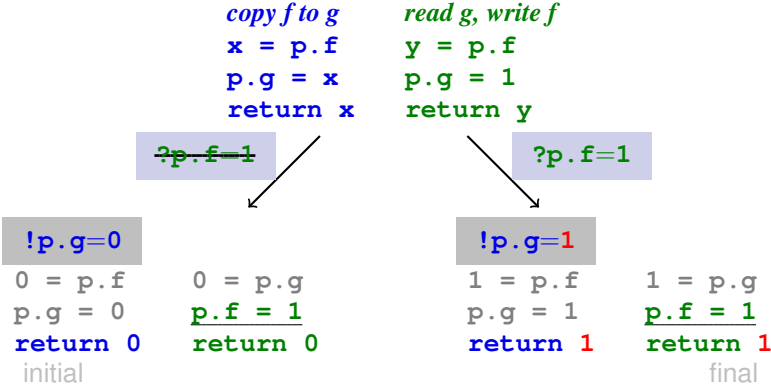
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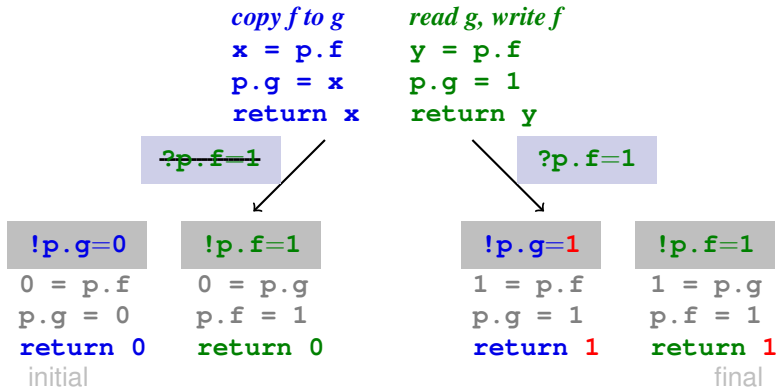
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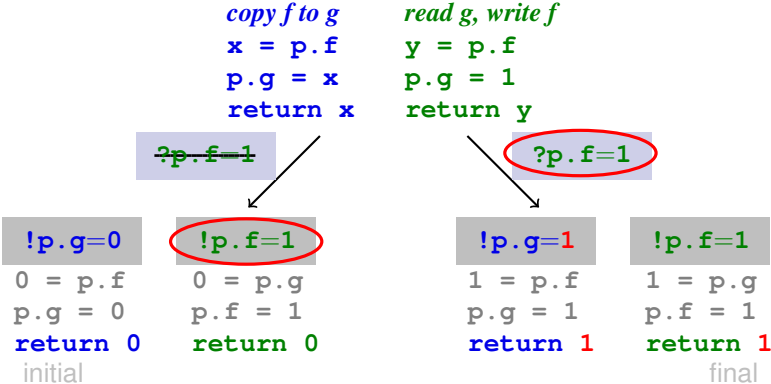
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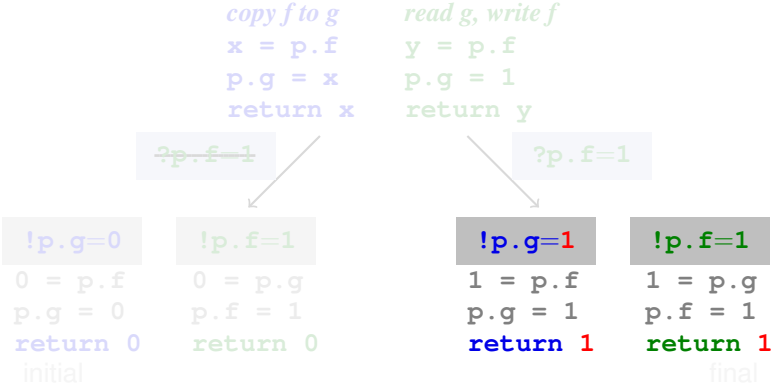
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Initiality necessary: Program B

- Initial branch must **justify** speculation
- Otherwise, execution is stuck

copy f to g
x = p.f
p.g = x
return x

copy g to f
y = p.g
p.f = y
return y

~~?p.f=1~~

?p.f=1

x = p.f
p.g = x
return x
initial

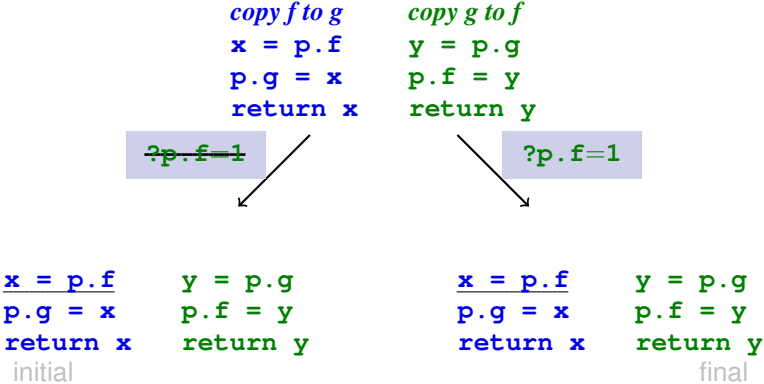
y = p.g
p.f = y
return y

x = p.f
p.g = x
return x

y = p.g
p.f = y
return y
final

Initiality necessary: Program B

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Initiality necessary: Program B

- Initial branch must **justify** speculation
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```

copy f to g
x = p.f
p.g = x
return x
    
```

```

copy g to f
y = p.g
p.f = y
return y
    
```

~~?p.f=1~~

?p.f=1

```

0 = p.f
p.g = 0
return 0
initial
    
```

```

y = p.g
p.f = y
return y
    
```

```

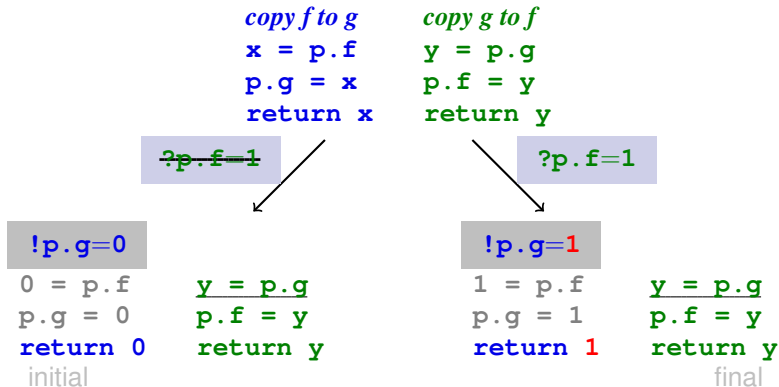
1 = p.f
p.g = 1
return 1
    
```

```

y = p.g
p.f = y
return y
final
    
```

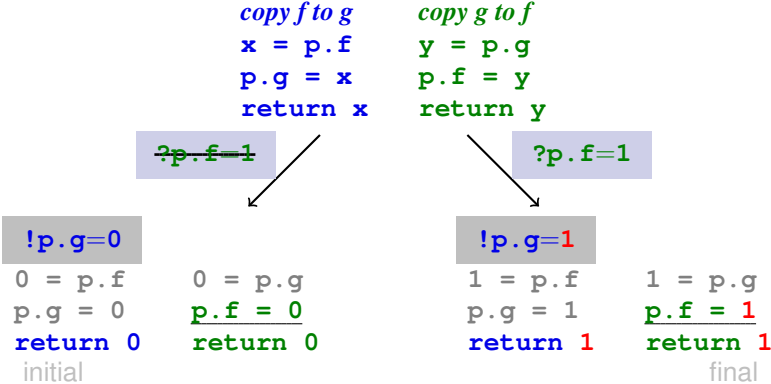
Initiality necessary: Program B

- Initial branch must **justify** speculation
- Otherwise, execution is stuck



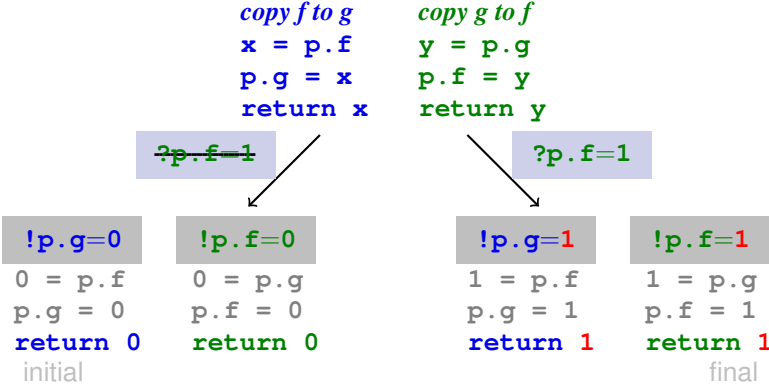
Initiality necessary: Program B

- Initial branch must **justify** speculation
- Otherwise, execution is stuck



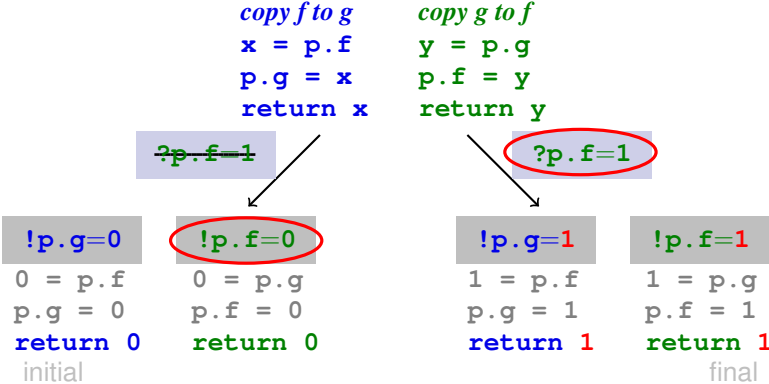
Initiality necessary: Program B

- Initial branch must **justify** speculation
- Otherwise, execution is stuck



Initiality necessary: Program B

- Initial branch must **justify** speculation
- Otherwise, execution is stuck



Initiality necessary: Program B

- Initial branch must **justify** speculation
- Otherwise, execution is stuck

```

copy f to g
x = p.f
p.g = x
return x

copy g to f
y = p.g
p.f = y
return y
    
```

~~?p.f=1~~

?p.f=1

STUCK!

!p.g=0

```

0 = p.f
p.g = 0
return 0
    
```

initial

!p.f=0

```

0 = p.g
p.f = 0
return 0
    
```

!p.g=1

```

1 = p.f
p.g = 1
return 1
    
```

!p.f=1

```

1 = p.g
p.f = 1
return 1
    
```

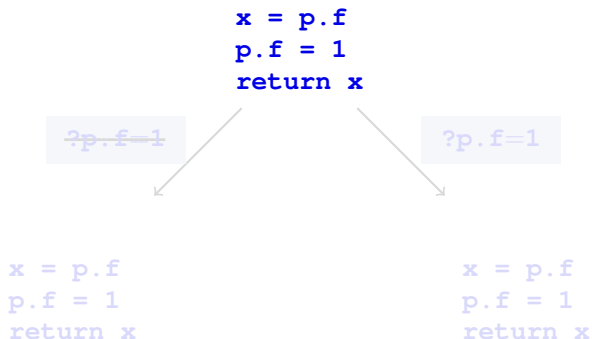
final

Avoiding thin air reads

- ✓ Initiality
- ? No self-justification
- ? Consistency
- ? Timeliness

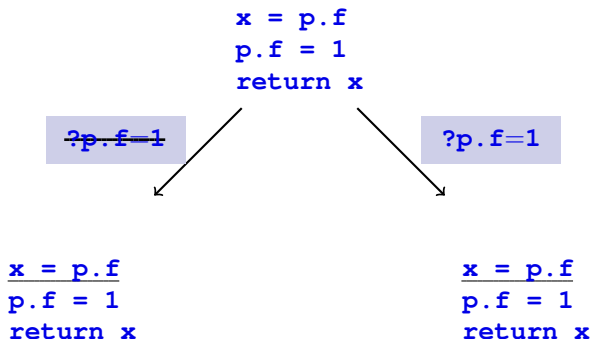
Self justification: a degenerate case

- Impossible in SC execution: **return 1**
- This execution prevented by definition of visibility:
Thread can not see self-speculation



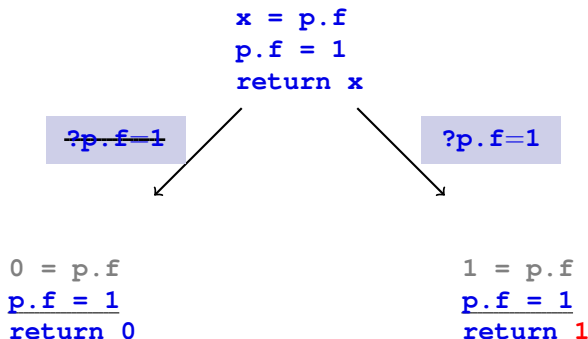
Self justification: a degenerate case

- Impossible in SC execution: **return 1**
- This execution prevented by definition of visibility:
Thread can not see self-speculation



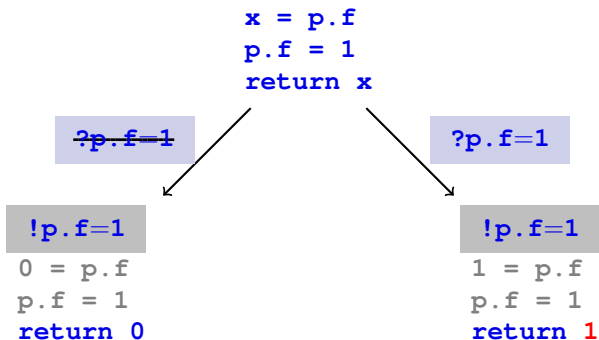
Self justification: a degenerate case

- Impossible in SC execution: **return 1**
- This execution prevented by definition of visibility:
Thread can not see self-speculation



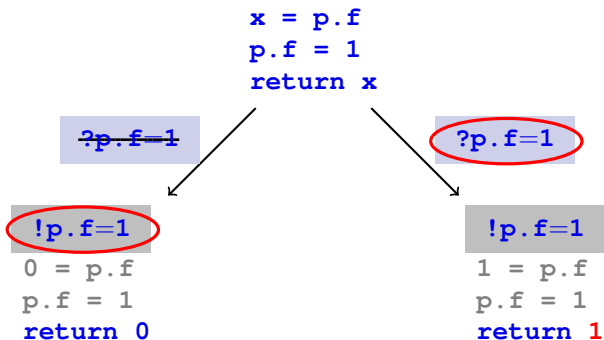
Self justification: a degenerate case

- Impossible in SC execution: **return 1**
- This execution prevented by definition of visibility:
Thread can not see self-speculation



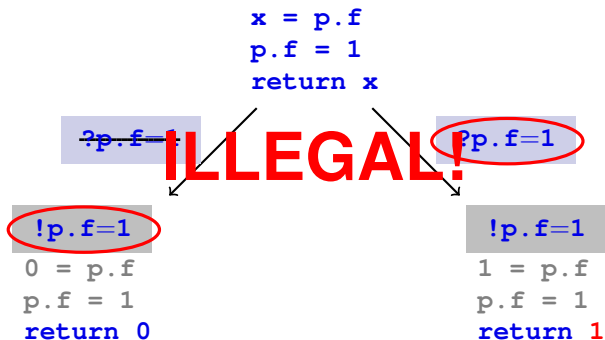
Self justification: a degenerate case

- Impossible in SC execution: **return 1**
- This execution prevented by definition of visibility:
Thread can not see self-speculation



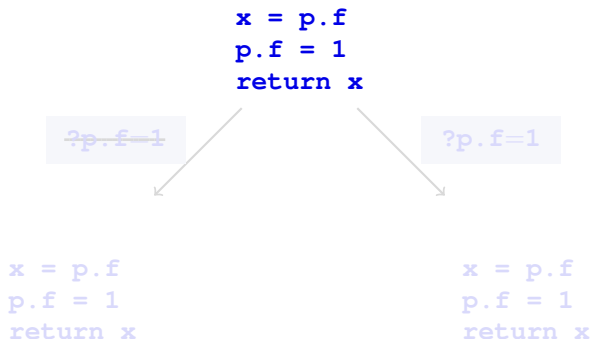
Self justification: a degenerate case

- Impossible in SC execution: `return 1`
- This execution prevented by definition of visibility:
Thread can not see self-speculation



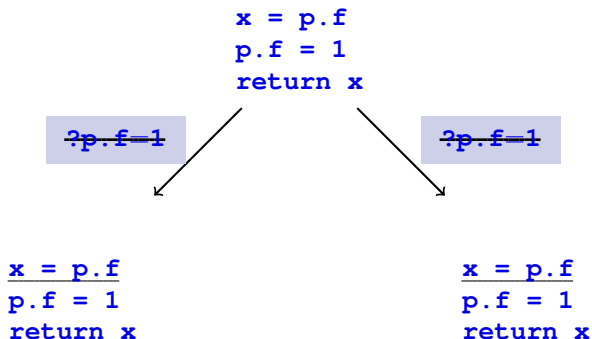
Self justification: a degenerate case

- Impossible in SC execution: **return 1**
- This execution prevented by definition of visibility:
Thread can not see self-speculation



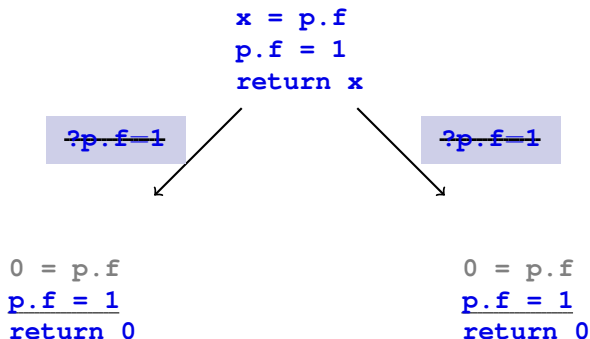
Self justification: a degenerate case

- Impossible in SC execution: `return 1`
- This execution prevented by definition of visibility:
Thread can not see self-speculation



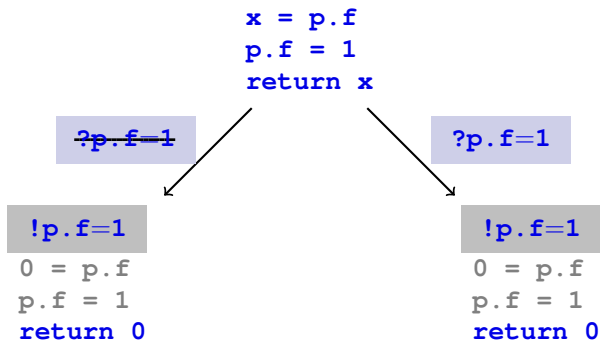
Self justification: a degenerate case

- Impossible in SC execution: `return 1`
- This execution prevented by definition of visibility:
Thread can not see self-speculation



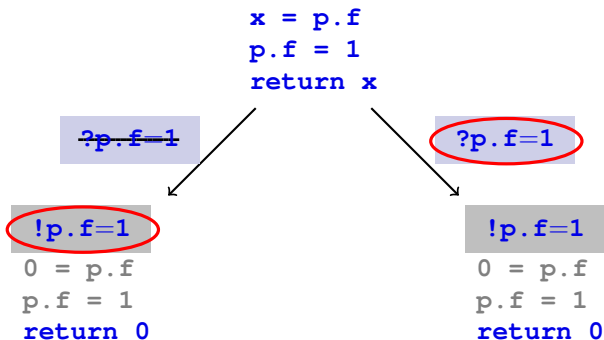
Self justification: a degenerate case

- Impossible in SC execution: `return 1`
- This execution prevented by definition of visibility:
Thread can not see self-speculation



Self justification: a degenerate case

- Impossible in SC execution: **return 1**
- This execution prevented by definition of visibility:
Thread can not see self-speculation



Controlling speculation: consistency

```
k.acquire()   k.acquire()   k.acquire()   k.acquire()
x = p.f      x = p.f      x = p.f      x = p.f
if(x==0)     if(x==0)     p.g = x      y = p.g
  p.f = 1    p.f = 2    k.release()  k.release()
k.release()  k.release()                return(x,y)
```

- Impossible in SC execution: `return(1,2)`
- Possible in final branch with speculation `?p.f=2`
- Initial branch can produce justifying write
- Inconsistent use of locks between speculation and justifying write

Initial branch

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>?p.f=2</code>
<code>⌋k</code>	<code>!p.g=0</code>	<code>k⌋</code>
<code>⌋k</code>	<code>!p.f=2</code>	<code>k⌋</code>
<code>⌋k</code>	<code>k⌋</code>	
<code>⌋k</code>	<code>k⌋</code>	

Final branch

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>?p.f=2</code>
<code>⌋k</code>	<code>!p.g=2</code>	<code>k⌋</code>
<code>⌋k</code>	<code>!p.f=1</code>	<code>k⌋</code>
<code>⌋k</code>	<code>k⌋</code>	
<code>⌋k</code>	<code>k⌋</code>	

Controlling speculation: consistency

```
k.acquire()   k.acquire()   k.acquire()   k.acquire()
  x = p.f     x = p.f     x = p.f     x = p.f
  if(x==0)   if(x==0)   p.g = x     y = p.g
    p.f = 1   p.f = 2   k.release() k.release()
k.release()  k.release()                return(x,y)
```

- Impossible in SC execution: `return(1,2)`
- Possible in final branch with speculation `?p.f=2`
- Initial branch can produce justifying write
- Inconsistent use of locks between speculation and justifying write

Initial branch

!p.f=0	!p.g=0	?p.f=2
↪k	!p.g=0	k↪
↪k	!p.f=2	k↪
↪k	k↪	
↪k	k↪	

Final branch

!p.f=0	!p.g=0	?p.f=2
↪k	!p.g=2	k↪
↪k	!p.f=1	k↪
↪k	k↪	
↪k	k↪	

Controlling speculation: consistency

```

k.acquire()   k.acquire()   k.acquire()   k.acquire()
  x = p.f     x = p.f     x = p.f     x = p.f
  if(x==0)    if(x==0)    p.g = x     y = p.g
    p.f = 1    p.f = 2    k.release() k.release()
k.release()   k.release()

```

- Impossible in SC execution: `return(1, 2)`
- Possible in final branch with speculation `?p.f=2`
- Initial branch can produce justifying write
- Inconsistent use of locks between speculation and justifying write

Initial branch

!p.f=0	!p.g=0	?p.f=2
↪k	!p.g=0	k↪
↪k	!p.f=2	k↪
↪k	k↪	
↪k	k↪	

Final branch

!p.f=0	!p.g=0	?p.f=2
↪k	!p.g=2	k↪
↪k	!p.f=1	k↪
↪k	k↪	
↪k	k↪	

Controlling speculation: consistency

```

k.acquire()   k.acquire()   k.acquire()   k.acquire()
  x = p.f     x = p.f     x = p.f     x = p.f
  if(x==0)    if(x==0)    p.g = x     y = p.g
    p.f = 1   p.f = 2   k.release() k.release()
k.release()   k.release()

```

- Impossible in SC execution: `return(1, 2)`
- Possible in final branch with speculation `?p.f=2`
- Initial branch can produce justifying write
- Inconsistent use of locks between speculation and justifying write

Initial branch

!p.f=0	!p.g=0
↪k	!p.g=0
↪k	!p.f=2
↪k	k↪
↪k	k↪

Final branch

!p.f=0	!p.g=0	?p.f=2
↪k	!p.g=2	k↪
↪k	!p.f=1	k↪
↪k	k↪	
↪k	k↪	

STUCK!

Controlling speculation: timeliness

```
k.acquire()   k.acquire()   k.acquire()
  x = p.f     x = p.f
  p.f = x+1   p.f = x+1
  p.g = 1     p.g = 2           y = p.g
k.release()   k.release()   k.release()
return x      return x      return y
```

- Impossible SC: **return 0; return 1; return 1;**
- Possible in final branch with speculation `?p.g=1`
- Initial and final branches produce same actions
- Speculation used to introduce data race in final branch

!p.f=0	!p.g=0			
$\frown k$!p.f=1	!p.g=1	$k \frown$?p.g=1
$\frown k$!p.f=2	!p.g=2	$k \frown$	
$\frown k$	$k \frown$			

Controlling speculation: timeliness

```
k.acquire()   k.acquire()   k.acquire()
  x = p.f     x = p.f     y = p.g
  p.f = x+1   p.f = x+1
  p.g = 1     p.g = 2
k.release()   k.release()   k.release()
return x      return x      return y
```

- Impossible SC: `return 0; return 1; return 1;`
- Possible in final branch with speculation `?p.g=1`
- Initial and final branches produce same actions
- Speculation used to introduce data race in final branch

<code>!p.f=0</code>	<code>!p.g=0</code>			
<code>↪k</code>	<code>!p.f=1</code>	<code>!p.g=1</code>	<code>k↪</code>	<code>?p.g=1</code>
<code>↪k</code>	<code>!p.f=2</code>	<code>!p.g=2</code>	<code>k↪</code>	
<code>↪k</code>	<code>k↪</code>			

Controlling speculation: timeliness

<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code> x = p.f</code>	<code> x = p.f</code>	
<code> p.f = x+1</code>	<code> p.f = x+1</code>	
<code> p.g = 1</code>	<code> p.g = 2</code>	<code> y = p.g</code>
<code>k.release()</code>	<code>k.release()</code>	<code>k.release()</code>
<code>return x</code>	<code>return x</code>	<code>return y</code>

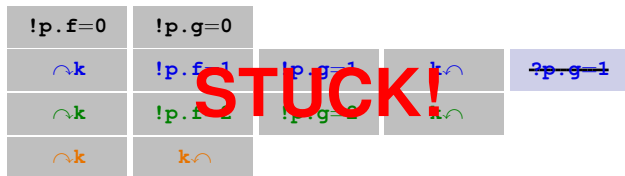
- Impossible SC: `return 0; return 1; return 1;`
- Possible in final branch with speculation `?p.g=1`
- Initial and final branches produce same actions
- Speculation used to introduce data race in final branch

<code>!p.f=0</code>	<code>!p.g=0</code>			
<code>↪k</code>	<code>!p.f=1</code>	<code>!p.g=1</code>	<code>k↪</code>	<code>?p.g=1</code>
<code>↪k</code>	<code>!p.f=2</code>	<code>!p.g=2</code>	<code>k↪</code>	
<code>↪k</code>	<code>k↪</code>			

Controlling speculation: timeliness

<pre> k.acquire() x = p.f p.f = x+1 p.g = 1 k.release() return x </pre>	<pre> k.acquire() x = p.f p.f = x+1 p.g = 2 k.release() return x </pre>	<pre> k.acquire() y = p.g k.release() return y </pre>
---	---	---

- Impossible SC: `return 0; return 1; return 1;`
- Possible in final branch with speculation `?p.g=1`
- Initial and final branches produce same actions
- Speculation used to introduce data race in final branch



1 Background

- Sequential Consistency
- Data Race Free Model
- Java Memory Model

2 Speculative semantics

- Empirical and speculative actions
- Desirable executions allowed
- Undesirable executions prevented

3 Summary of results

- Relation to Java Memory Model
- Simulation precongruence

Relation to JMM

Theorem

DRF program \Rightarrow SC execution

Theorem

Lockless program \Rightarrow every JMM execution allowed

Simulation

- Simulation defined in paper
- Precongruence
- Useful

$$\begin{aligned} & \mathbf{p.f=1; p.g=1; } \succsim \mathbf{p.g=1; p.f=1;} \\ & \mathbf{p.f=1; k.acquire(); } \succsim \mathbf{k.acquire(); p.f=1;} \\ & \mathbf{x=p.f; y=p.f; } M \succsim \mathbf{x=p.f; } M\{\mathbf{x/y}\} \end{aligned}$$

Summary

- New model based on **speculation**

Data Races	Locks	New vs. JMM
X	-	=
-	X	>
✓	✓	≠

- Simulation precongruence
- Better behaved:
Validates redundant-read-elimination, roach-motel, etc
- Thank you

Summary

- New model based on **speculation**

Data Races	Locks	New vs. JMM
X	-	=
-	X	>
✓	✓	≠

- Simulation precongruence
- Better behaved:
Validates redundant-read-elimination, roach-motel, etc
- Thank you

Appendix

The rest of the slides animate the execution of a few examples.

Controlling speculation: consistency

<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code> x = p.f</code>	<code> x = p.f</code>	<code> x = p.f</code>	<code> x = p.f</code>
<code> if(x==0)</code>	<code> if(x==0)</code>	<code> p.g = x</code>	<code> y = p.g</code>
<code> p.f = 1</code>	<code> p.f = 2</code>	<code>k.release()</code>	<code>k.release()</code>
<code>k.release()</code>	<code>k.release()</code>		<code>return(x,y)</code>

<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code> x = p.f</code>	<code> x = p.f</code>	<code> x = p.f</code>	<code> x = p.f</code>
<code> if(x==0)</code>	<code> if(x==0)</code>	<code> p.g = x</code>	<code> y = p.g</code>
<code> p.f = 1</code>	<code> p.f = 2</code>	<code>k.release()</code>	<code>k.release()</code>
<code>k.release()</code>	<code>k.release()</code>		<code>return(x,y)</code>

Controlling speculation: consistency

!p.f=0

!p.g=0

~~?p.f=2~~

```
k.acquire()  k.acquire()
x = p.f      x = p.f
if(x==0)     if(x==0)
  p.f = 1    p.f = 2
k.release()  k.release()
```

```
k.acquire()  k.acquire()
x = p.f      x = p.f
p.g = x      y = p.g
k.release()  k.release()
return(x,y)
```

!p.f=0

!p.g=0

?p.f=2

```
k.acquire()  k.acquire()
x = p.f      x = p.f
if(x==0)     if(x==0)
  p.f = 1    p.f = 2
k.release()  k.release()
```

```
k.acquire()  k.acquire()
x = p.f      x = p.f
p.g = x      y = p.g
k.release()  k.release()
return(x,y)
```

Controlling speculation: consistency

!p.f=0

!p.g=0

~~?p.f=2~~

\cap k

<pre>k.acquire() x = p.f if(x==0) p.f = 1 k.release()</pre>	<pre>k.acquire() x = p.f if(x==0) p.f = 2 k.release()</pre>	<pre>k.acquire() <u>x = p.f</u> <u>p.g = x</u> k.release()</pre>	<pre>k.acquire() x = p.f y = p.g k.release() return(x,y)</pre>
---	---	--	--

!p.f=0

!p.g=0

?p.f=2

\cap k

<pre>k.acquire() x = p.f if(x==0) p.f = 1 k.release()</pre>	<pre>k.acquire() x = p.f if(x==0) p.f = 2 k.release()</pre>	<pre>k.acquire() <u>x = p.f</u> <u>p.g = x</u> k.release()</pre>	<pre>k.acquire() x = p.f y = p.g k.release() return(x,y)</pre>
---	---	--	--

Controlling speculation: consistency

!p.f=0

!p.g=0

~~?p.f=2~~

\cap k

```
k.acquire()
x = p.f
if(x==0)
  p.f = 1
k.release()
```

```
k.acquire()
x = p.f
if(x==0)
  p.f = 2
k.release()
```

```
k.acquire()
0 = p.f
p.g = 0
k.release()
```

```
k.acquire()
x = p.f
y = p.g
k.release()
return(x,y)
```

!p.f=0

!p.g=0

?p.f=2

\cap k

```
k.acquire()
x = p.f
if(x==0)
  p.f = 1
k.release()
```

```
k.acquire()
x = p.f
if(x==0)
  p.f = 2
k.release()
```

```
k.acquire()
2 = p.f
p.g = 2
k.release()
```

```
k.acquire()
x = p.f
y = p.g
k.release()
return(x,y)
```

Controlling speculation: consistency

`!p.f=0`

`!p.g=0`

~~`?p.f=2`~~

`⊃k`

`!p.g=0`

<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code>x = p.f</code>	<code>x = p.f</code>	<code>0 = p.f</code>	<code>x = p.f</code>
<code>if(x==0)</code>	<code>if(x==0)</code>	<code>p.g = 0</code>	<code>y = p.g</code>
<code> p.f = 1</code>	<code> p.f = 2</code>	<u><code>k.release()</code></u>	<code>k.release()</code>
<code>k.release()</code>	<code>k.release()</code>		<code>return(x,y)</code>

`!p.f=0`

`!p.g=0`

`?p.f=2`

`⊃k`

`!p.g=2`

<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code>x = p.f</code>	<code>x = p.f</code>	<code>2 = p.f</code>	<code>x = p.f</code>
<code>if(x==0)</code>	<code>if(x==0)</code>	<code>p.g = 2</code>	<code>y = p.g</code>
<code> p.f = 1</code>	<code> p.f = 2</code>	<u><code>k.release()</code></u>	<code>k.release()</code>
<code>k.release()</code>	<code>k.release()</code>		<code>return(x,y)</code>

Controlling speculation: consistency

`!p.f=0`

`!p.g=0`

~~`?p.f=2`~~

`∩k`

`!p.g=0`

`k∩`

`k.acquire()`

`x = p.f`

`if(x==0)`

`p.f = 1`

`k.release()`

`k.acquire()`

`x = p.f`

`if(x==0)`

`p.f = 2`

`k.release()`

`k.acquire()`

`0 = p.f`

`p.g = 0`

`k.release()`

`k.acquire()`

`x = p.f`

`y = p.g`

`k.release()`

`return(x,y)`

`!p.f=0`

`!p.g=0`

`?p.f=2`

`∩k`

`!p.g=2`

`k∩`

`k.acquire()`

`x = p.f`

`if(x==0)`

`p.f = 1`

`k.release()`

`k.acquire()`

`x = p.f`

`if(x==0)`

`p.f = 2`

`k.release()`

`k.acquire()`

`2 = p.f`

`p.g = 2`

`k.release()`

`k.acquire()`

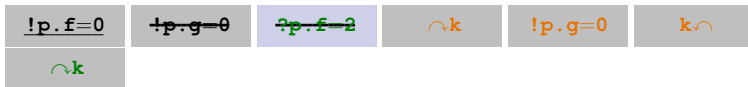
`x = p.f`

`y = p.g`

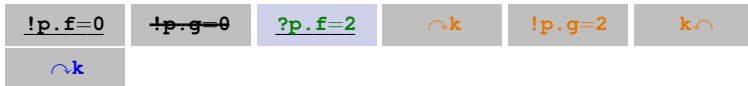
`k.release()`

`return(x,y)`

Controlling speculation: consistency

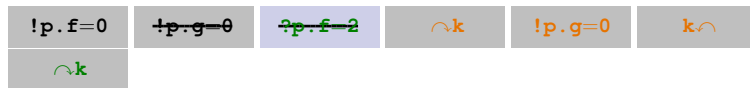


<pre>k.acquire() x = p.f if(x==0) p.f = 1 k.release()</pre>	<pre>k.acquire() <u>x = p.f</u> if(x==0) p.f = 2 k.release()</pre>	<pre>k.acquire() 0 = p.f p.g = 0 k.release()</pre>	<pre>k.acquire() x = p.f y = p.g k.release() return(x,y)</pre>
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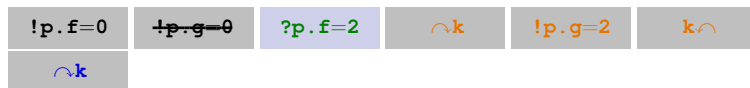


<pre>k.acquire() <u>x = p.f</u> if(x==0) p.f = 1 k.release()</pre>	<pre>k.acquire() x = p.f if(x==0) p.f = 2 k.release()</pre>	<pre>k.acquire() 2 = p.f p.g = 2 k.release()</pre>	<pre>k.acquire() x = p.f y = p.g k.release() return(x,y)</pre>
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Controlling speculation: consistency

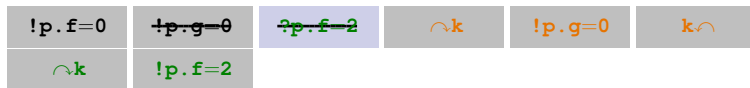


<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code>x = p.f</code>	<code>0 = p.f</code>	<code>0 = p.f</code>	<code>x = p.f</code>
<code>if(x==0)</code>	<code>if(0==0)</code>	<code>p.g = 0</code>	<code>y = p.g</code>
<code> p.f = 1</code>	<code> <u>p.f = 2</u></code>	<code>k.release()</code>	<code>k.release()</code>
<code>k.release()</code>	<code>k.release()</code>		<code>return(x,y)</code>

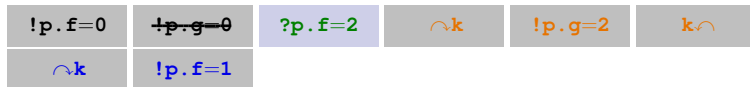


<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code>0 = p.f</code>	<code>x = p.f</code>	<code>2 = p.f</code>	<code>x = p.f</code>
<code>if(0==0)</code>	<code>if(x==0)</code>	<code>p.g = 2</code>	<code>y = p.g</code>
<code> <u>p.f = 1</u></code>	<code> p.f = 2</code>	<code>k.release()</code>	<code>k.release()</code>
<code>k.release()</code>	<code>k.release()</code>		<code>return(x,y)</code>

Controlling speculation: consistency



<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code>x = p.f</code>	<code>0 = p.f</code>	<code>0 = p.f</code>	<code>x = p.f</code>
<code>if(x==0)</code>	<code>if(0==0)</code>	<code>p.g = 0</code>	<code>y = p.g</code>
<code> p.f = 1</code>	<code> p.f = 2</code>	<code>k.release()</code>	<code>k.release()</code>
<code>k.release()</code>	<u><code>k.release()</code></u>		<code>return(x,y)</code>



<code>k.acquire()</code>	<u><code>k.acquire()</code></u>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code>0 = p.f</code>	<code>x = p.f</code>	<code>2 = p.f</code>	<code>x = p.f</code>
<code>if(0==0)</code>	<code>if(x==0)</code>	<code>p.g = 2</code>	<code>y = p.g</code>
<code> p.f = 1</code>	<code> p.f = 2</code>	<code>k.release()</code>	<code>k.release()</code>
<u><code>k.release()</code></u>	<code>k.release()</code>		<code>return(x,y)</code>

Controlling speculation: consistency

$\neg p.f=0$	$\neg p.g=0$	$?p.f=2$	$\neg k$	$\neg p.g=0$	$k \wedge$
$\neg k$	$\neg p.f=2$	$k \wedge$			

<u>$k.acquire()$</u> $x = p.f$ $if(x==0)$ $p.f = 1$ $k.release()$	$k.acquire()$ $0 = p.f$ $if(0==0)$ $p.f = 2$ $k.release()$	$k.acquire()$ $0 = p.f$ $p.g = 0$ $k.release()$	$k.acquire()$ $x = p.f$ $y = p.g$ $k.release()$ $return(x, y)$
--	--	--	--

$\neg p.f=0$	$\neg p.g=0$	$?p.f=2$	$\neg k$	$\neg p.g=2$	$k \wedge$
$\neg k$	$\neg p.f=1$	$k \wedge$			

$k.acquire()$ $0 = p.f$ $if(0==0)$ $p.f = 1$ $k.release()$	<u>$k.acquire()$</u> $x = p.f$ $if(x==0)$ $p.f = 2$ $k.release()$	$k.acquire()$ $2 = p.f$ $p.g = 2$ $k.release()$	$k.acquire()$ $x = p.f$ $y = p.g$ $k.release()$ $return(x, y)$
--	--	--	--

Controlling speculation: consistency

!p.f=0	!p.g=0	?p.f=2	$\cap k$!p.g=0	$k \cap$
$\cap k$	<u>!p.f=2</u>	$k \cap$	$\cap k$		
<pre>k.acquire() <u>x = p.f</u> if(x==0) p.f = 1 k.release()</pre>	<pre>k.acquire() 0 = p.f if(0==0) p.f = 2 k.release()</pre>	<pre>k.acquire() 0 = p.f p.g = 0 k.release()</pre>	<pre>k.acquire() x = p.f y = p.g k.release() return(x,y)</pre>		

!p.f=0	!p.g=0	?p.f=2	$\cap k$!p.g=2	$k \cap$
$\cap k$	<u>!p.f=1</u>	$k \cap$	$\cap k$		
<pre>k.acquire() 0 = p.f if(0==0) p.f = 1 k.release()</pre>	<pre>k.acquire() <u>x = p.f</u> if(x==0) p.f = 2 k.release()</pre>	<pre>k.acquire() 2 = p.f p.g = 2 k.release()</pre>	<pre>k.acquire() x = p.f y = p.g k.release() return(x,y)</pre>		

Controlling speculation: consistency

!p.f=0	!p.g=0	?p.f=2	$\cap k$!p.g=0	$k \cap$
$\cap k$!p.f=2	$k \cap$	$\cap k$		
<pre>k.acquire() 2 = p.f if(2==0) p.f = 1 <u>k.release()</u></pre>	<pre>k.acquire() 0 = p.f if(0==0) p.f = 2 k.release()</pre>	<pre>k.acquire() 0 = p.f p.g = 0 k.release()</pre>	<pre>k.acquire() x = p.f y = p.g k.release() return(x,y)</pre>		

!p.f=0	!p.g=0	?p.f=2	$\cap k$!p.g=2	$k \cap$
$\cap k$!p.f=1	$k \cap$	$\cap k$		
<pre>k.acquire() 0 = p.f if(0==0) p.f = 1 k.release()</pre>	<pre>k.acquire() 1 = p.f if(1==0) p.f = 2 <u>k.release()</u></pre>	<pre>k.acquire() 2 = p.f p.g = 2 k.release()</pre>	<pre>k.acquire() x = p.f y = p.g k.release() return(x,y)</pre>		

Controlling speculation: consistency

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>?p.f=2</code>	<code>⊃k</code>	<code>!p.g=0</code>	<code>k⊃</code>
<code>⊃k</code>	<code>!p.f=2</code>	<code>k⊃</code>	<code>⊃k</code>	<code>k⊃</code>	
<code>k.acquire()</code> <code>2 = p.f</code> <code>if(2==0)</code> <code> p.f = 1</code> <code>k.release()</code>	<code>k.acquire()</code> <code>0 = p.f</code> <code>if(0==0)</code> <code> p.f = 2</code> <code>k.release()</code>	<code>k.acquire()</code> <code>0 = p.f</code> <code>if(0==0)</code> <code> p.g = 0</code> <code>k.release()</code>	<code>k.acquire()</code> <code>0 = p.f</code> <code>p.g = 0</code> <code>k.release()</code>	<u><code>k.acquire()</code></u> <u><code>x = p.f</code></u> <u><code>y = p.g</code></u> <u><code>k.release()</code></u> <u><code>return(x,y)</code></u>	

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>?p.f=2</code>	<code>⊃k</code>	<code>!p.g=2</code>	<code>k⊃</code>
<code>⊃k</code>	<code>!p.f=1</code>	<code>k⊃</code>	<code>⊃k</code>	<code>k⊃</code>	
<code>k.acquire()</code> <code>0 = p.f</code> <code>if(0==0)</code> <code> p.f = 1</code> <code>k.release()</code>	<code>k.acquire()</code> <code>1 = p.f</code> <code>if(1==0)</code> <code> p.f = 2</code> <code>k.release()</code>	<code>k.acquire()</code> <code>2 = p.f</code> <code>if(2==0)</code> <code> p.g = 2</code> <code>k.release()</code>	<code>k.acquire()</code> <code>2 = p.f</code> <code>p.g = 2</code> <code>k.release()</code>	<u><code>k.acquire()</code></u> <u><code>x = p.f</code></u> <u><code>y = p.g</code></u> <u><code>k.release()</code></u> <u><code>return(x,y)</code></u>	

Controlling speculation: consistency

!p.f=0	!p.g=0	?p.f=2	$\cap k$!p.g=0	$k \cap$
$\cap k$	<u>!p.f=2</u>	$k \cap$	$\cap k$	$k \cap$	$\cap k$

<pre>k.acquire() 2 = p.f if(2==0) p.f = 1 k.release()</pre>	<pre>k.acquire() 0 = p.f if(0==0) p.f = 2 k.release()</pre>	<pre>k.acquire() 0 = p.f p.g = 0 k.release()</pre>	<pre>k.acquire() <u>x = p.f</u> y = p.g k.release() return(x,y)</pre>
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!p.f=0	!p.g=0	?p.f=2	$\cap k$!p.g=2	$k \cap$
$\cap k$	<u>!p.f=1</u>	$k \cap$	$\cap k$	$k \cap$	$\cap k$

<pre>k.acquire() 0 = p.f if(0==0) p.f = 1 k.release()</pre>	<pre>k.acquire() 1 = p.f if(1==0) p.f = 2 k.release()</pre>	<pre>k.acquire() 2 = p.f p.g = 2 k.release()</pre>	<pre>k.acquire() <u>x = p.f</u> y = p.g k.release() return(x,y)</pre>
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Controlling speculation: consistency

!p.f=0	!p.g=0	?p.f=2	$\cap k$	<u>!p.g=0</u>	$k \cap$
$\cap k$!p.f=2	$k \cap$	$\cap k$	$k \cap$	$\cap k$

<pre>k.acquire() 2 = p.f if(2==0) p.f = 1 k.release()</pre>	<pre>k.acquire() 0 = p.f if(0==0) p.f = 2 k.release()</pre>	<pre>k.acquire() 0 = p.f p.g = 0 k.release()</pre>	<pre>k.acquire() 2 = p.f <u>y = p.g</u> k.release() return(2, y)</pre>
---	---	--	--

!p.f=0	!p.g=0	?p.f=2	$\cap k$	<u>!p.g=2</u>	$k \cap$
$\cap k$!p.f=1	$k \cap$	$\cap k$	$k \cap$	$\cap k$

<pre>k.acquire() 0 = p.f if(0==0) p.f = 1 k.release()</pre>	<pre>k.acquire() 1 = p.f if(1==0) p.f = 2 k.release()</pre>	<pre>k.acquire() 2 = p.f p.g = 2 k.release()</pre>	<pre>k.acquire() 1 = p.f <u>y = p.g</u> k.release() return(1, y)</pre>
---	---	--	--

Controlling speculation: consistency

!p.f=0	!p.g=0	?p.f=2	$\cap k$!p.g=0	$k \cap$
$\cap k$!p.f=2	$k \cap$	$\cap k$	$k \cap$	$\cap k$
k.acquire() 2 = p.f if(2==0) p.f = 1 k.release()	k.acquire() 0 = p.f if(0==0) p.f = 2 k.release()	k.acquire() 0 = p.f if(0==0) p.f = 2 k.release()	k.acquire() 0 = p.f p.g = 0 k.release()	k.acquire() 2 = p.f 0 = p.g <u>k.release()</u> return(2,0)	

!p.f=0	!p.g=0	?p.f=2	$\cap k$!p.g=2	$k \cap$
$\cap k$!p.f=1	$k \cap$	$\cap k$	$k \cap$	$\cap k$
k.acquire() 0 = p.f if(0==0) p.f = 1 k.release()	k.acquire() 1 = p.f if(1==0) p.f = 2 k.release()	k.acquire() 1 = p.f if(1==0) p.f = 2 k.release()	k.acquire() 2 = p.f p.g = 2 k.release()	k.acquire() 1 = p.f 2 = p.g <u>k.release()</u> return(1,2)	

Controlling speculation: consistency

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>?p.f=2</code>	<code>⊃k</code>	<code>!p.g=0</code>	<code>k⊃</code>	
<code>⊃k</code>	<code>!p.f=2</code>	<code>k⊃</code>	<code>⊃k</code>	<code>k⊃</code>	<code>⊃k</code>	<code>k⊃</code>
<code>k.acquire()</code> <code>2 = p.f</code> <code>if(2==0)</code> <code> p.f = 1</code> <code>k.release()</code>	<code>k.acquire()</code> <code>0 = p.f</code> <code>if(0==0)</code> <code> p.f = 2</code> <code>k.release()</code>	<code>k.acquire()</code> <code>0 = p.f</code> <code>if(0==0)</code> <code> p.g = 0</code> <code>k.release()</code>	<code>k.acquire()</code> <code>0 = p.f</code> <code>p.g = 0</code> <code>k.release()</code>	<code>k.acquire()</code> <code>2 = p.f</code> <code>0 = p.g</code> <code>k.release()</code>	<code>return(2, 0)</code>	

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>?p.f=2</code>	<code>⊃k</code>	<code>!p.g=2</code>	<code>k⊃</code>	
<code>⊃k</code>	<code>!p.f=1</code>	<code>k⊃</code>	<code>⊃k</code>	<code>k⊃</code>	<code>⊃k</code>	<code>k⊃</code>
<code>k.acquire()</code> <code>0 = p.f</code> <code>if(0==0)</code> <code> p.f = 1</code> <code>k.release()</code>	<code>k.acquire()</code> <code>1 = p.f</code> <code>if(1==0)</code> <code> p.f = 2</code> <code>k.release()</code>	<code>k.acquire()</code> <code>2 = p.f</code> <code>p.g = 2</code> <code>k.release()</code>	<code>k.acquire()</code> <code>2 = p.f</code> <code>p.g = 2</code> <code>k.release()</code>	<code>k.acquire()</code> <code>1 = p.f</code> <code>2 = p.g</code> <code>k.release()</code>	<code>return(1, 2)</code>	

Controlling speculation: timeliness

<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code>x = p.f</code>	<code>x = p.f</code>	
<code>p.f = x+1</code>	<code>p.f = x+1</code>	
<code>p.g = 1</code>	<code>p.g = 2</code>	<code>y = p.g</code>
<code>k.release()</code>	<code>k.release()</code>	<code>k.release()</code>
<code>return x</code>	<code>return x</code>	<code>return y</code>

<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code>x = p.f</code>	<code>x = p.f</code>	
<code>p.f = x+1</code>	<code>p.f = x+1</code>	
<code>p.g = 1</code>	<code>p.g = 2</code>	<code>y = p.g</code>
<code>k.release()</code>	<code>k.release()</code>	<code>k.release()</code>
<code>return x</code>	<code>return x</code>	<code>return y</code>

Controlling speculation: timeliness

!p.f=0

!p.g=0

<u>k.acquire()</u>	k.acquire()	k.acquire()
x = p.f	x = p.f	
p.f = x+1	p.f = x+1	
p.g = 1	p.g = 2	y = p.g
k.release()	k.release()	k.release()
return x	return x	return y

!p.f=0

!p.g=0

<u>k.acquire()</u>	k.acquire()	k.acquire()
x = p.f	x = p.f	
p.f = x+1	p.f = x+1	
p.g = 1	p.g = 2	y = p.g
k.release()	k.release()	k.release()
return x	return x	return y

Controlling speculation: timeliness

!p.f=0

!p.g=0

\rightsquigarrow k

k.acquire() <u>x = p.f</u> p.f = x+1 p.g = 1 k.release() return x	k.acquire() x = p.f p.f = x+1 p.g = 2 k.release() return x	k.acquire() y = p.g k.release() return y
--	---	---

!p.f=0

!p.g=0

\rightsquigarrow k

k.acquire() <u>x = p.f</u> p.f = x+1 p.g = 1 k.release() return x	k.acquire() x = p.f p.f = x+1 p.g = 2 k.release() return x	k.acquire() y = p.g k.release() return y
--	---	---

Controlling speculation: timeliness

!p.f=0

!p.g=0

⋈k

k.acquire() 0 = p.f <u>p.f = 0+1</u> p.g = 1 k.release() return 0	k.acquire() x = p.f p.f = x+1 p.g = 2 k.release() return x	k.acquire() y = p.g k.release() return y
--	---	---

!p.f=0

!p.g=0

⋈k

k.acquire() 0 = p.f <u>p.f = 0+1</u> p.g = 1 k.release() return 0	k.acquire() x = p.f p.f = x+1 p.g = 2 k.release() return x	k.acquire() y = p.g k.release() return y
--	---	---

Controlling speculation: timeliness

!p.f=0

!p.g=0

\leadsto k

!p.f=1

k.acquire() 0 = p.f p.f = 0+1 <u>p.g = 1</u> k.release() return 0	k.acquire() x = p.f p.f = x+1 p.g = 2 k.release() return x	k.acquire() y = p.g k.release() return y
--	---	---

!p.f=0

!p.g=0

\leadsto k

!p.f=1

k.acquire() 0 = p.f p.f = 0+1 <u>p.g = 1</u> k.release() return 0	k.acquire() x = p.f p.f = x+1 p.g = 2 k.release() return x	k.acquire() y = p.g k.release() return y
--	---	---

Controlling speculation: timeliness

!p.f=0

!p.g=0

\curvearrowright k

!p.f=1

!p.g=1

```
k.acquire()
  0 = p.f
  p.f = 0+1
  p.g = 1
k.release()
return 0
```

```
k.acquire()
  x = p.f
  p.f = x+1
  p.g = 2
k.release()
return x
```

```
k.acquire()
  y = p.g
k.release()
return y
```

!p.f=0

!p.g=0

\curvearrowright k

!p.f=1

!p.g=1

```
k.acquire()
  0 = p.f
  p.f = 0+1
  p.g = 1
k.release()
return 0
```

```
k.acquire()
  x = p.f
  p.f = x+1
  p.g = 2
k.release()
return x
```

```
k.acquire()
  y = p.g
k.release()
return y
```

Controlling speculation: timeliness

!p.f=0

!p.g=0

\curvearrowright k

!p.f=1

!p.g=1

k \curvearrowright

k.acquire() 0 = p.f p.f = 0+1 p.g = 1 k.release() return 0	k.acquire() x = p.f p.f = x+1 p.g = 2 k.release() return x	k.acquire() y = p.g k.release() return y
---	---	---

!p.f=0

!p.g=0

\curvearrowright k

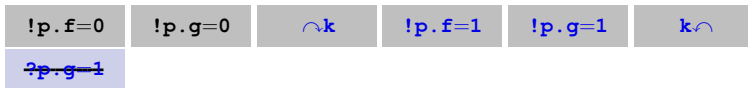
!p.f=1

!p.g=1

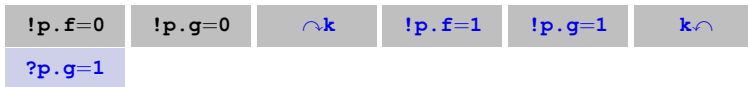
k \curvearrowright

k.acquire() 0 = p.f p.f = 0+1 p.g = 1 k.release() return 0	k.acquire() x = p.f p.f = x+1 p.g = 2 k.release() return x	k.acquire() y = p.g k.release() return y
---	---	---

Controlling speculation: timeliness

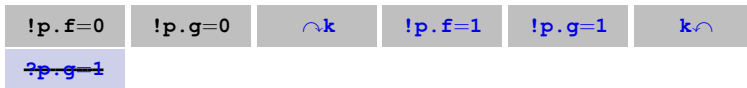


<code>k.acquire()</code> <code>0 = p.f</code> <code>p.f = 0+1</code> <code>p.g = 1</code> <code>k.release()</code> <code>return 0</code>	<code>k.acquire()</code> <code>x = p.f</code> <code>p.f = x+1</code> <code>p.g = 2</code> <code>k.release()</code> <code>return x</code>	<code>k.acquire()</code> <code>y = p.g</code> <code>k.release()</code> <code>return y</code>
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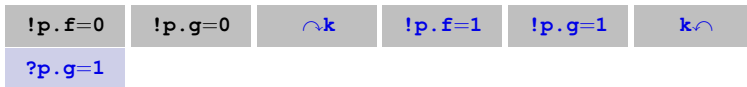


<code>k.acquire()</code> <code>0 = p.f</code> <code>p.f = 0+1</code> <code>p.g = 1</code> <code>k.release()</code> <code>return 0</code>	<code>k.acquire()</code> <code>x = p.f</code> <code>p.f = x+1</code> <code>p.g = 2</code> <code>k.release()</code> <code>return x</code>	<code>k.acquire()</code> <code>y = p.g</code> <code>k.release()</code> <code>return y</code>
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Controlling speculation: timeliness

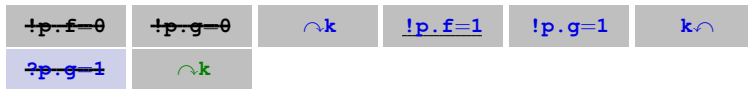


<code>k.acquire()</code>	<code><u>k.acquire()</u></code>	<code>k.acquire()</code>
<code> 0 = p.f</code>	<code> x = p.f</code>	
<code> p.f = 0+1</code>	<code> p.f = x+1</code>	
<code> p.g = 1</code>	<code> p.g = 2</code>	<code> y = p.g</code>
<code>k.release()</code>	<code>k.release()</code>	<code>k.release()</code>
<code>return 0</code>	<code>return x</code>	<code>return y</code>

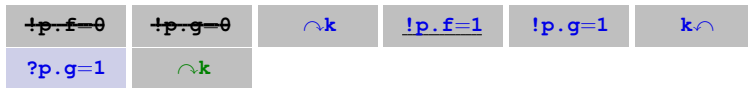


<code>k.acquire()</code>	<code><u>k.acquire()</u></code>	<code>k.acquire()</code>
<code> 0 = p.f</code>	<code> x = p.f</code>	
<code> p.f = 0+1</code>	<code> p.f = x+1</code>	
<code> p.g = 1</code>	<code> p.g = 2</code>	<code> y = p.g</code>
<code>k.release()</code>	<code>k.release()</code>	<code>k.release()</code>
<code>return 0</code>	<code>return x</code>	<code>return y</code>

Controlling speculation: timeliness

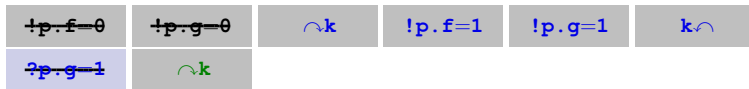


<pre>k.acquire() 0 = p.f p.f = 0+1 p.g = 1 k.release() return 0</pre>	<pre>k.acquire() <u>x = p.f</u> p.f = x+1 p.g = 2 k.release() return x</pre>	<pre>k.acquire() y = p.g k.release() return y</pre>
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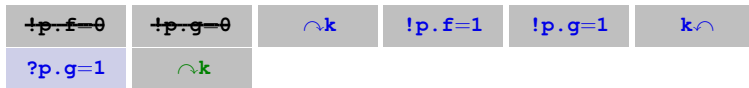


<pre>k.acquire() 0 = p.f p.f = 0+1 p.g = 1 k.release() return 0</pre>	<pre>k.acquire() <u>x = p.f</u> p.f = x+1 p.g = 2 k.release() return x</pre>	<pre>k.acquire() y = p.g k.release() return y</pre>
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Controlling speculation: timeliness



<pre>k.acquire() 0 = p.f p.f = 0+1 p.g = 1 k.release() return 0</pre>	<pre>k.acquire() 1 = p.f <u>p.f = 1+1</u> p.g = 2 k.release() return 1</pre>	<pre>k.acquire() y = p.g k.release() return y</pre>
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<pre>k.acquire() 0 = p.f p.f = 0+1 p.g = 1 k.release() return 0</pre>	<pre>k.acquire() 1 = p.f <u>p.f = 1+1</u> p.g = 2 k.release() return 1</pre>	<pre>k.acquire() y = p.g k.release() return y</pre>
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Controlling speculation: timeliness

!p.f=0	!p.g=0	$\neg k$!p.f=1	!p.g=1	$k \wedge$
?p.g=1	$\neg k$!p.f=2			

<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code>0 = p.f</code>	<code>1 = p.f</code>	
<code>p.f = 0+1</code>	<code>p.f = 1+1</code>	
<code>p.g = 1</code>	<u><code>p.g = 2</code></u>	<code>y = p.g</code>
<code>k.release()</code>	<u><code>k.release()</code></u>	<code>k.release()</code>
<code>return 0</code>	<code>return 1</code>	<code>return y</code>

!p.f=0	!p.g=0	$\neg k$!p.f=1	!p.g=1	$k \wedge$
?p.g=1	$\neg k$!p.f=2			

<code>k.acquire()</code>	<code>k.acquire()</code>	<code>k.acquire()</code>
<code>0 = p.f</code>	<code>1 = p.f</code>	
<code>p.f = 0+1</code>	<code>p.f = 1+1</code>	
<code>p.g = 1</code>	<u><code>p.g = 2</code></u>	<code>y = p.g</code>
<code>k.release()</code>	<u><code>k.release()</code></u>	<code>k.release()</code>
<code>return 0</code>	<code>return 1</code>	<code>return y</code>

Controlling speculation: timeliness

!p.f=0	!p.g=0	\curvearrowright k	!p.f=1	!p.g=1	k \curvearrowright
?p.g=1	\curvearrowright k	!p.f=2	!p.g=2		

k.acquire() 0 = p.f p.f = 0+1 p.g = 1 k.release() return 0	k.acquire() 1 = p.f p.f = 1+1 p.g = 2 <u>k.release()</u> return 1	k.acquire() y = p.g k.release() return y
---	--	---

!p.f=0	!p.g=0	\curvearrowright k	!p.f=1	!p.g=1	k \curvearrowright
?p.g=1	\curvearrowright k	!p.f=2	!p.g=2		

k.acquire() 0 = p.f p.f = 0+1 p.g = 1 k.release() return 0	k.acquire() 1 = p.f p.f = 1+1 p.g = 2 <u>k.release()</u> return 1	k.acquire() y = p.g k.release() return y
---	--	---

Controlling speculation: timeliness

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>⌊k</code>	<code>!p.f=1</code>	<code>!p.g=1</code>	<code>k⌋</code>
<code>?p.g=1</code>	<code>⌊k</code>	<code>!p.f=2</code>	<code>!p.g=2</code>	<code>k⌋</code>	

<code>k.acquire()</code>	<code>k.acquire()</code>	<u><code>k.acquire()</code></u>
<code>0 = p.f</code>	<code>1 = p.f</code>	
<code>p.f = 0+1</code>	<code>p.f = 1+1</code>	
<code>p.g = 1</code>	<code>p.g = 2</code>	<code>y = p.g</code>
<code>k.release()</code>	<code>k.release()</code>	<u><code>k.release()</code></u>
<code>return 0</code>	<code>return 1</code>	<code>return y</code>

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>⌊k</code>	<code>!p.f=1</code>	<code>!p.g=1</code>	<code>k⌋</code>
<code>?p.g=1</code>	<code>⌊k</code>	<code>!p.f=2</code>	<code>!p.g=2</code>	<code>k⌋</code>	

<code>k.acquire()</code>	<code>k.acquire()</code>	<u><code>k.acquire()</code></u>
<code>0 = p.f</code>	<code>1 = p.f</code>	
<code>p.f = 0+1</code>	<code>p.f = 1+1</code>	
<code>p.g = 1</code>	<code>p.g = 2</code>	<code>y = p.g</code>
<code>k.release()</code>	<code>k.release()</code>	<u><code>k.release()</code></u>
<code>return 0</code>	<code>return 1</code>	<code>return y</code>

Controlling speculation: timeliness

!p.f=0	!p.g=0	$\curvearrowright k$!p.f=1	!p.g=1	$k \curvearrowright$
<u>?p.g=1</u>	$\curvearrowright k$!p.f=2	<u>!p.g=2</u>	$k \curvearrowright$	$\curvearrowright k$

<pre>k.acquire() 0 = p.f p.f = 0+1 p.g = 1 k.release() return 0</pre>	<pre>k.acquire() 1 = p.f p.f = 1+1 p.g = 2 k.release() return 1</pre>	<pre>k.acquire() <u>y = p.g</u> k.release() return y</pre>
---	---	--

!p.f=0	!p.g=0	$\curvearrowright k$!p.f=1	!p.g=1	$k \curvearrowright$
<u>?p.g=1</u>	$\curvearrowright k$!p.f=2	<u>!p.g=2</u>	$k \curvearrowright$	$\curvearrowright k$

<pre>k.acquire() 0 = p.f p.f = 0+1 p.g = 1 k.release() return 0</pre>	<pre>k.acquire() 1 = p.f p.f = 1+1 p.g = 2 k.release() return 1</pre>	<pre>k.acquire() <u>y = p.g</u> k.release() return y</pre>
---	---	--

Controlling speculation: timeliness

!p.f=0	!p.g=0	$\curvearrowright k$!p.f=1	!p.g=1	$k \curvearrowright$
?p.g=1	$\curvearrowright k$!p.f=2	!p.g=2	$k \curvearrowright$	$\curvearrowright k$

<pre>k.acquire() 0 = p.f p.f = 0+1 p.g = 1 k.release() return 0</pre>	<pre>k.acquire() 1 = p.f p.f = 1+1 p.g = 2 k.release() return 1</pre>	<pre>k.acquire() 2 = p.g <u>k.release()</u> return 2</pre>
---	---	--

!p.f=0	!p.g=0	$\curvearrowright k$!p.f=1	!p.g=1	$k \curvearrowright$
?p.g=1	$\curvearrowright k$!p.f=2	!p.g=2	$k \curvearrowright$	$\curvearrowright k$

<pre>k.acquire() 0 = p.f p.f = 0+1 p.g = 1 k.release() return 0</pre>	<pre>k.acquire() 1 = p.f p.f = 1+1 p.g = 2 k.release() return 1</pre>	<pre>k.acquire() 1 = p.g <u>k.release()</u> return 1</pre>
---	---	--

Controlling speculation: timeliness

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>↪k</code>	<code>!p.f=1</code>	<code>!p.g=1</code>	<code>k↪</code>	
<code>?p.g=1</code>	<code>↪k</code>	<code>!p.f=2</code>	<code>!p.g=2</code>	<code>k↪</code>	<code>↪k</code>	<code>k↪</code>

```

k.acquire()    k.acquire()    k.acquire()
 0 = p.f       1 = p.f
 p.f = 0+1     p.f = 1+1
 p.g = 1
k.release()    k.release()    k.release()
return 0       return 1       return 2
  
```

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>↪k</code>	<code>!p.f=1</code>	<code>!p.g=1</code>	<code>k↪</code>	
<code>?p.g=1</code>	<code>↪k</code>	<code>!p.f=2</code>	<code>!p.g=2</code>	<code>k↪</code>	<code>↪k</code>	<code>k↪</code>

```

k.acquire()    k.acquire()    k.acquire()
 0 = p.f       1 = p.f
 p.f = 0+1     p.f = 1+1
 p.g = 1
k.release()    k.release()    k.release()
return 0       return 1       return 1
  
```

Roach Motel

```
k.acquire()   k.acquire()   x = p.f      y = p.g
  p.f = 2     p.f = 1     k.acquire()  p.h = y
k.release()   k.release()  z = p.h      return y
              if(x==2) p.g = 1
              else    p.g = z
k.release()
return(x, z)
```

```
k.acquire()   k.acquire()   x = p.f      y = p.g
  p.f = 2     p.f = 1     k.acquire()  p.h = y
k.release()   k.release()  z = p.h      return y
              if(x==2) p.g = 1
              else    p.g = z
k.release()
return(x, z)
```

Roach Motel

!p.f=0

!p.g=0

!p.h=0

```
k.acquire()   k.acquire()  
  p.f = 2      p.f = 1  
k.release()    k.release()
```

```
x = p.f  
k.acquire()  
  z = p.h  
  if(x==2) p.g = 1  
  else    p.g = z  
k.release()  
return(x, z)
```

```
y = p.g  
p.h = y  
return y
```

!p.f=0

!p.g=0

!p.h=0

```
k.acquire()   k.acquire()  
  p.f = 2      p.f = 1  
k.release()    k.release()
```

```
x = p.f  
k.acquire()  
  z = p.h  
  if(x==2) p.g = 1  
  else    p.g = z  
k.release()  
return(x, z)
```

```
y = p.g  
p.h = y  
return y
```

Roach Motel

`!p.f=0`

`!p.g=0`

`!p.h=0`

`⊃k`

`k.acquire()`

`p.f = 2`

`k.release()`

`k.acquire()`

`p.f = 1`

`k.release()`

`x = p.f`

`k.acquire()`

`z = p.h`

`if(x==2) p.g = 1`

`else p.g = z`

`k.release()`

`return(x, z)`

`y = p.g`

`p.h = y`

`return y`

`!p.f=0`

`!p.g=0`

`!p.h=0`

`⊃k`

`k.acquire()`

`p.f = 2`

`k.release()`

`k.acquire()`

`p.f = 1`

`k.release()`

`x = p.f`

`k.acquire()`

`z = p.h`

`if(x==2) p.g = 1`

`else p.g = z`

`k.release()`

`return(x, z)`

`y = p.g`

`p.h = y`

`return y`

Roach Motel

!p.f=0

!p.g=0

!p.h=0

⊃k

!p.f=2

```
k.acquire()   k.acquire()
  p.f = 2     p.f = 1
k.release() k.release()
```

```
x = p.f
k.acquire()
  z = p.h
  if(x==2) p.g = 1
  else    p.g = z
k.release()
return(x, z)
```

```
y = p.g
p.h = y
return y
```

!p.f=0

!p.g=0

!p.h=0

⊃k

!p.f=2

```
k.acquire()   k.acquire()
  p.f = 2     p.f = 1
k.release() k.release()
```

```
x = p.f
k.acquire()
  z = p.h
  if(x==2) p.g = 1
  else    p.g = z
k.release()
return(x, z)
```

```
y = p.g
p.h = y
return y
```

Roach Motel

!p.f=0

!p.g=0

!p.h=0

\cap k

!p.f=2

k \cap

```
k.acquire()   k.acquire()
  p.f = 2      p.f = 1
k.release()   k.release()
```

```
x = p.f
k.acquire()
  z = p.h
  if(x==2) p.g = 1
  else    p.g = z
k.release()
return(x, z)
```

```
y = p.g
p.h = y
return y
```

!p.f=0

!p.g=0

!p.h=0

\cap k

!p.f=2

k \cap

```
k.acquire()   k.acquire()
  p.f = 2      p.f = 1
k.release()   k.release()
```

```
x = p.f
k.acquire()
  z = p.h
  if(x==2) p.g = 1
  else    p.g = z
k.release()
return(x, z)
```

```
y = p.g
p.h = y
return y
```

Roach Motel

`!p.f=0`

`!p.g=0`

`!p.h=0`

`⊃k`

`!p.f=2`

`k⊃`

```
k.acquire()   k.acquire()
  p.f = 2     p.f = 1
k.release()   k.release()
```

```
2 = p.f
k.acquire()
  z = p.h
  if(2==2) p.g = 1
  else    p.g = z
k.release()
return(2, z)
```

```
y = p.g
p.h = y
return y
```

~~`!p.f=0`~~

`!p.g=0`

`!p.h=0`

`⊃k`

`!p.f=2`

`k⊃`

`⊃k`

```
k.acquire()   k.acquire()
  p.f = 2     p.f = 1
k.release()   k.release()
```

```
x = p.f
k.acquire()
  z = p.h
  if(x==2) p.g = 1
  else    p.g = z
k.release()
return(x, z)
```

```
y = p.g
p.h = y
return y
```

Roach Motel

~~!p.f=0~~

!p.g=0

!p.h=0

$\curvearrowright k$

!p.f=2

$k \curvearrowright$

$\curvearrowright k$

```
k.acquire()   k.acquire()
  p.f = 2      p.f = 1
k.release()   k.release()
```

```
2 = p.f
k.acquire()
  z = p.h
  if(2==2) p.g = 1
  else     p.g = z
k.release()
return(2, z)
```

```
y = p.g
p.h = y
return y
```

~~!p.f=0~~

!p.g=0

!p.h=0

$\curvearrowright k$

!p.f=2

$k \curvearrowright$

$\curvearrowright k$

!p.f=1

```
k.acquire()   k.acquire()
  p.f = 2      p.f = 1
k.release()   k.release()
```

```
x = p.f
k.acquire()
  z = p.h
  if(x==2) p.g = 1
  else     p.g = z
k.release()
return(x, z)
```

```
y = p.g
p.h = y
return y
```

Roach Motel



```
k.acquire()    k.acquire()    2 = p.f
  p.f = 2      p.f = 1      k.acquire()
k.release()    k.release()    z = p.h
                if(2==2) p.g = 1
                else    p.g = z
                k.release()
                return(2, z)

y = p.g
p.h = y
return y
```



```
k.acquire()    k.acquire()    x = p.f
  p.f = 2      p.f = 1      k.acquire()
k.release()    k.release()    z = p.h
                if(x==2) p.g = 1
                else    p.g = z
                k.release()
                return(x, z)

y = p.g
p.h = y
return y
```

Roach Motel



```
k.acquire()   k.acquire()
  p.f = 2     p.f = 1
k.release()   k.release()
```

```
2 = p.f
k.acquire()
  z = p.h
  if(2==2) p.g = 1
  else    p.g = z
k.release()
return(2, z)
```

```
y = p.g
p.h = y
return y
```



```
k.acquire()   k.acquire()
  p.f = 2     p.f = 1
k.release()   k.release()
```

```
1 = p.f
k.acquire()
  z = p.h
  if(1==2) p.g = 1
  else    p.g = z
k.release()
return(1, z)
```

```
y = p.g
p.h = y
return y
```

Roach Motel



```
k.acquire()    k.acquire()
  p.f = 2      p.f = 1
k.release()    k.release()
```

```
2 = p.f
k.acquire()
  z = p.h
  if(2==2) p.g = 1
  else    p.g = z
k.release()
return(2, z)
```

```
y = p.g
p.h = y
return y
```



```
k.acquire()    k.acquire()
  p.f = 2      p.f = 1
k.release()    k.release()
```

```
1 = p.f
k.acquire()
  z = p.h
  if(1==2) p.g = 1
  else    p.g = z
k.release()
return(1, z)
```

```
y = p.g
p.h = y
return y
```

Roach Motel

!p.f=0	!p.g=0	!p.h=0	$\curvearrowright k$!p.f=2	$k \curvearrowright$	$\curvearrowright k$
!p.f=1	$k \curvearrowright$	$\curvearrowright k$?p.h=1			
<pre>k.acquire() p.f = 2 k.release()</pre>	<pre>k.acquire() p.f = 1 k.release()</pre>	<pre>2 = p.f k.acquire() z = p.h if(2==2) p.g = 1 else p.g = z k.release() return(2, z)</pre>			<pre>y = p.g p.h = y return y</pre>	

!p.f=0	!p.g=0	!p.h=0	$\curvearrowright k$!p.f=2	$k \curvearrowright$	$\curvearrowright k$
!p.f=1	$k \curvearrowright$	$\curvearrowright k$?p.h=1			
<pre>k.acquire() p.f = 2 k.release()</pre>	<pre>k.acquire() p.f = 1 k.release()</pre>	<pre>1 = p.f k.acquire() z = p.h if(1==2) p.g = 1 else p.g = z k.release() return(1, z)</pre>			<pre>y = p.g p.h = y return y</pre>	

Roach Motel

!p.f=0	!p.g=0	<u>!p.h=0</u>	$\cap k$!p.f=2	$k \cap$	$\cap k$
!p.f=1	$k \cap$	$\cap k$	<u>?p.h=1</u>			
<pre>k.acquire() p.f = 2 k.release()</pre>	<pre>k.acquire() p.f = 1 k.release()</pre>	<pre>2 = p.f k.acquire() <u>z = p.h</u> if(2==2) p.g = 1 else p.g = z k.release() return(2, z)</pre>	<pre>y = p.g p.h = y return y</pre>			

!p.f=0	!p.g=0	<u>!p.h=0</u>	$\cap k$!p.f=2	$k \cap$	$\cap k$
!p.f=1	$k \cap$	$\cap k$	<u>?p.h=1</u>			
<pre>k.acquire() p.f = 2 k.release()</pre>	<pre>k.acquire() p.f = 1 k.release()</pre>	<pre>1 = p.f k.acquire() <u>z = p.h</u> if(1==2) p.g = 1 else p.g = z k.release() return(1, z)</pre>	<pre>y = p.g p.h = y return y</pre>			

Roach Motel

!p.f=0	!p.g=0	!p.h=0	$\curvearrowright k$!p.f=2	$k \curvearrowright$	$\curvearrowright k$
!p.f=1	$k \curvearrowright$	$\curvearrowright k$?p.h=1			
<pre>k.acquire() p.f = 2 k.release()</pre>	<pre>k.acquire() p.f = 1 k.release()</pre>	<pre>2 = p.f k.acquire() 0 = p.h if(2==2) <u>p.g = 1</u> else p.g = 0 k.release() return(2, 0)</pre>			<pre>y = p.g p.h = y return y</pre>	

!p.f=0	!p.g=0	!p.h=0	$\curvearrowright k$!p.f=2	$k \curvearrowright$	$\curvearrowright k$
!p.f=1	$k \curvearrowright$	$\curvearrowright k$?p.h=1			
<pre>k.acquire() p.f = 2 k.release()</pre>	<pre>k.acquire() p.f = 1 k.release()</pre>	<pre>1 = p.f k.acquire() 1 = p.h if(1==2) p.g = 1 else <u>p.g = 1</u> k.release() return(1, 1)</pre>			<pre>y = p.g p.h = y return y</pre>	

Roach Motel

<code>!p.f=0</code>	<code><u>!p.g=0</u></code>	<code>!p.h=0</code>	<code>↪k</code>	<code>!p.f=2</code>	<code>k↪</code>	<code>↪k</code>
<code>!p.f=1</code>	<code>k↪</code>	<code>↪k</code>	<code>?p.h=1</code>	<code><u>!p.g=1</u></code>		
<code>k.acquire()</code> <code>p.f = 2</code> <code>k.release()</code>	<code>k.acquire()</code> <code>p.f = 1</code> <code>k.release()</code>	<code>2 = p.f</code> <code>k.acquire()</code> <code>0 = p.h</code> <code>if(2==2) p.g = 1</code> <code>else p.g = 0</code> <code>k.release()</code> <code>return(2, 0)</code>		<code>y = p.g</code> <code>p.h = y</code> <code>return y</code>		

<code>!p.f=0</code>	<code><u>!p.g=0</u></code>	<code>!p.h=0</code>	<code>↪k</code>	<code>!p.f=2</code>	<code>k↪</code>	<code>↪k</code>
<code>!p.f=1</code>	<code>k↪</code>	<code>↪k</code>	<code>?p.h=1</code>	<code><u>!p.g=1</u></code>		
<code>k.acquire()</code> <code>p.f = 2</code> <code>k.release()</code>	<code>k.acquire()</code> <code>p.f = 1</code> <code>k.release()</code>	<code>1 = p.f</code> <code>k.acquire()</code> <code>1 = p.h</code> <code>if(1==2) p.g = 1</code> <code>else p.g = 1</code> <code>k.release()</code> <code>return(1, 1)</code>		<code>y = p.g</code> <code>p.h = y</code> <code>return y</code>		

Roach Motel

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>!p.h=0</code>	<code>↪k</code>	<code>!p.f=2</code>	<code>k↪</code>	<code>↪k</code>
<code>!p.f=1</code>	<code>k↪</code>	<code>↪k</code>	<code>?p.h=1</code>	<code>!p.g=1</code>		
<code>k.acquire()</code> <code>p.f = 2</code> <code>k.release()</code>	<code>k.acquire()</code> <code>p.f = 1</code> <code>k.release()</code>	<code>2 = p.f</code> <code>k.acquire()</code> <code>0 = p.h</code> <code>if(2==2) p.g = 1</code> <code>else p.g = 0</code> <code>k.release()</code> <code>return(2, 0)</code>		<code>1 = p.g</code> <code><u>p.h = 1</u></code> <code>return 1</code>		

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>!p.h=0</code>	<code>↪k</code>	<code>!p.f=2</code>	<code>k↪</code>	<code>↪k</code>
<code>!p.f=1</code>	<code>k↪</code>	<code>↪k</code>	<code>?p.h=1</code>	<code>!p.g=1</code>		
<code>k.acquire()</code> <code>p.f = 2</code> <code>k.release()</code>	<code>k.acquire()</code> <code>p.f = 1</code> <code>k.release()</code>	<code>1 = p.f</code> <code>k.acquire()</code> <code>1 = p.h</code> <code>if(1==2) p.g = 1</code> <code>else p.g = 1</code> <code>k.release()</code> <code>return(1, 1)</code>		<code>1 = p.g</code> <code><u>p.h = 1</u></code> <code>return 1</code>		

Roach Motel

!p.f=0	!p.g=0	!p.h=0	$\curvearrowright k$!p.f=2	$k \curvearrowright$	$\curvearrowright k$
!p.f=1	$k \curvearrowright$	$\curvearrowright k$?p.h=1	!p.g=1	!p.h=1	
k.acquire() p.f = 2 k.release()	k.acquire() p.f = 1 k.release()	2 = p.f k.acquire() 0 = p.h if(2==2) p.g = 1 else p.g = 0 <u>k.release()</u> return(2, 0)		1 = p.g p.h = 1 return 1		

!p.f=0	!p.g=0	!p.h=0	$\curvearrowright k$!p.f=2	$k \curvearrowright$	$\curvearrowright k$
!p.f=1	$k \curvearrowright$	$\curvearrowright k$?p.h=1	!p.g=1	!p.h=1	
k.acquire() p.f = 2 k.release()	k.acquire() p.f = 1 k.release()	1 = p.f k.acquire() 1 = p.h if(1==2) p.g = 1 else p.g = 1 <u>k.release()</u> return(1, 1)		1 = p.g p.h = 1 return 1		

Roach Motel

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>!p.h=0</code>	<code>⊃k</code>	<code>!p.f=2</code>	<code>k⊃</code>	<code>⊃k</code>
<code>!p.f=1</code>	<code>k⊃</code>	<code>⊃k</code>	<code>?p.h=1</code>	<code>!p.g=1</code>	<code>!p.h=1</code>	<code>k⊃</code>

```
k.acquire()
p.f = 2
k.release()

k.acquire()
p.f = 1
k.release()
```

```
2 = p.f
k.acquire()
0 = p.h
if(2==2) p.g = 1
else     p.g = 0
k.release()
return(2, 0)
```

```
1 = p.g
p.h = 1
return 1
```

<code>!p.f=0</code>	<code>!p.g=0</code>	<code>!p.h=0</code>	<code>⊃k</code>	<code>!p.f=2</code>	<code>k⊃</code>	<code>⊃k</code>
<code>!p.f=1</code>	<code>k⊃</code>	<code>⊃k</code>	<code>?p.h=1</code>	<code>!p.g=1</code>	<code>!p.h=1</code>	<code>k⊃</code>

```
k.acquire()
p.f = 2
k.release()

k.acquire()
p.f = 1
k.release()
```

```
1 = p.f
k.acquire()
1 = p.h
if(1==2) p.g = 1
else     p.g = 1
k.release()
return(1, 1)
```

```
1 = p.g
p.h = 1
return 1
```